ECE2030b - HW-5 v. 2 Due Monday 10/21 during class. - ANSWERS
Problem 1. Using Finite State Machine techniques, design a circuit to:
Detect when sequential input $X$ delivers 3 logic 1 's in a row.
Do not detect overlapping sequences.
Example:
Input: 01101111011111001111110
Output: 00000010000100000010010
A. Draw a State Diagram showing all possible states and transitions.
B. Draw a logic table for the Next State bits ( Ni ) and the Output bit (Q), as a function of Present State bits (Pi) and Input bit (X).
C. Draw Karnaugh maps for the separate outputs, Ni and Q .
D. Draw a logic diagram showing the necessary registers and combinatorial logic blocks.
A. State Diagram (Meely)


State Diagram (Moore)


Check List: Does every state have exits defined for all inputs (0,1)?
B. Logic or Truth Tables:

| Meely |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Present | State | Input | Next | State | Output* |
| P1 | P0 | X | N1 | N0 | Q |
| 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 |
| 1 | 1 | 0 | 0 | 0 | 0 |
| 1 | 1 | 1 | 0 | 0 | 0 |

- For Meely Machine, output occurs while machine is in state 10 and $X=1$.

Moore Machine:

| Moore |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| Present | State | Input | Next | State | Output* |
| P1 | P0 | X | N1 | N0 | Q |
| 0 | 0 | 0 | 0 | 0 | 0 |
| 0 | 0 | 1 | 0 | 1 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 |
| 0 | 1 | 1 | 1 | 0 | 0 |
| 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{0}$ |
| 1 | 1 | 0 | 0 | 0 | $\mathbf{1}$ |
| 1 | 1 | 1 | 0 | $\mathbf{1}$ | $\mathbf{1}$ |

For Moore Machine, output occurs while machine is in state 11. Logic for $Q$ can be designed as a function of $N 1, N 0$
B. Karnaugh Maps for Moore Machine:

| $\mathrm{N} 1: \mathrm{X}$ | $\backslash$ | $\mathrm{P} 1, \mathrm{P} 0$ | 00 | 01 | 11 | 10 |
| ---: | ---: | ---: | :---: | :---: | :---: | :---: |
|  | 0 | 0 | 0 | 0 | 0 |  |
|  | 1 | 0 | 1 | 0 | 1 |  |

$\mathrm{N} 1=\mathrm{X}\left(\mathrm{P} 1^{\prime} \mathrm{P} 0+\mathrm{P} 1 \mathrm{P} 0^{\prime}\right)=\mathrm{X}(\mathrm{P} 1 \mathrm{XOR} \mathrm{P} 0$ )

| $\mathrm{N} 0: \mathrm{X}$ | $\backslash$ | $\mathrm{P} 1, \mathrm{P} 0$ | 00 | 01 | 11 | 10 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
|  | 0 | 0 | 0 | 0 | 0 |  |
|  | 1 | 1 | 0 | 1 | 1 |  |

$$
\mathrm{N} 0=\mathrm{X}\left(\mathrm{P} 0^{\prime}+\mathrm{P} 1\right)
$$

| Q: P1 \P0 | 0 | 1 |  |
| :--- | :--- | :--- | :--- |
|  | 0 | 0 | 0 |
|  | 1 | 0 | 1 |

Q = P0 P1 (note: for Moore Machine, $Q$ is function of present state (P1,P0).
D. Logic Diagram (Moore)


Solution as a Meely Machine is also acceptable.

